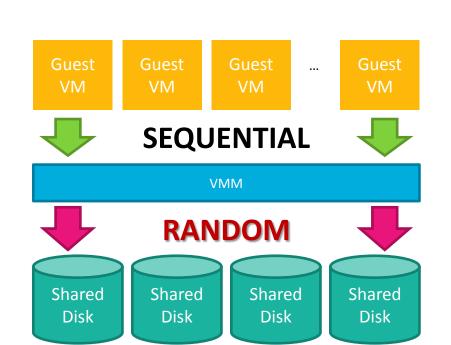
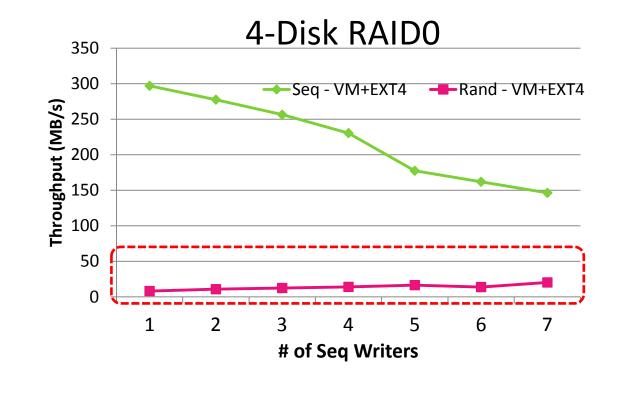
# Gecko: Contention-Oblivious Disk Arrays for Cloud Storage

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### Motivation

- Cloud/Virtualization accelerates consolidation of servers
  - Numbers of CPU cores and VMs increase per server
  - Storage is typically poorly virtualized





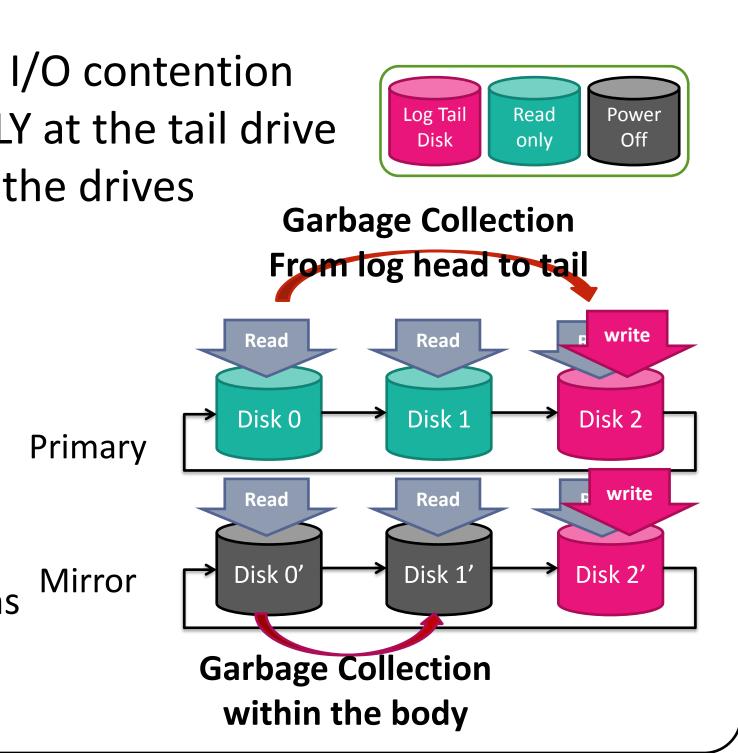
I/O contention destroy the performance

#### Challenges to I/O contention

- I/O contention to overcome
  - I/O contention moves the disk head and destroys performance
  - Write-write, write-read, read-read, write-GC, and read-GC
- RAID cannot preserve high throughput
  - IO performance varies depending on coexisting VMs
- Vulnerable to I/O Contention
- LFS only solves write-write contention
  - GC (Garbage Collection) interferes with logging
  - First class reads interfere with logging

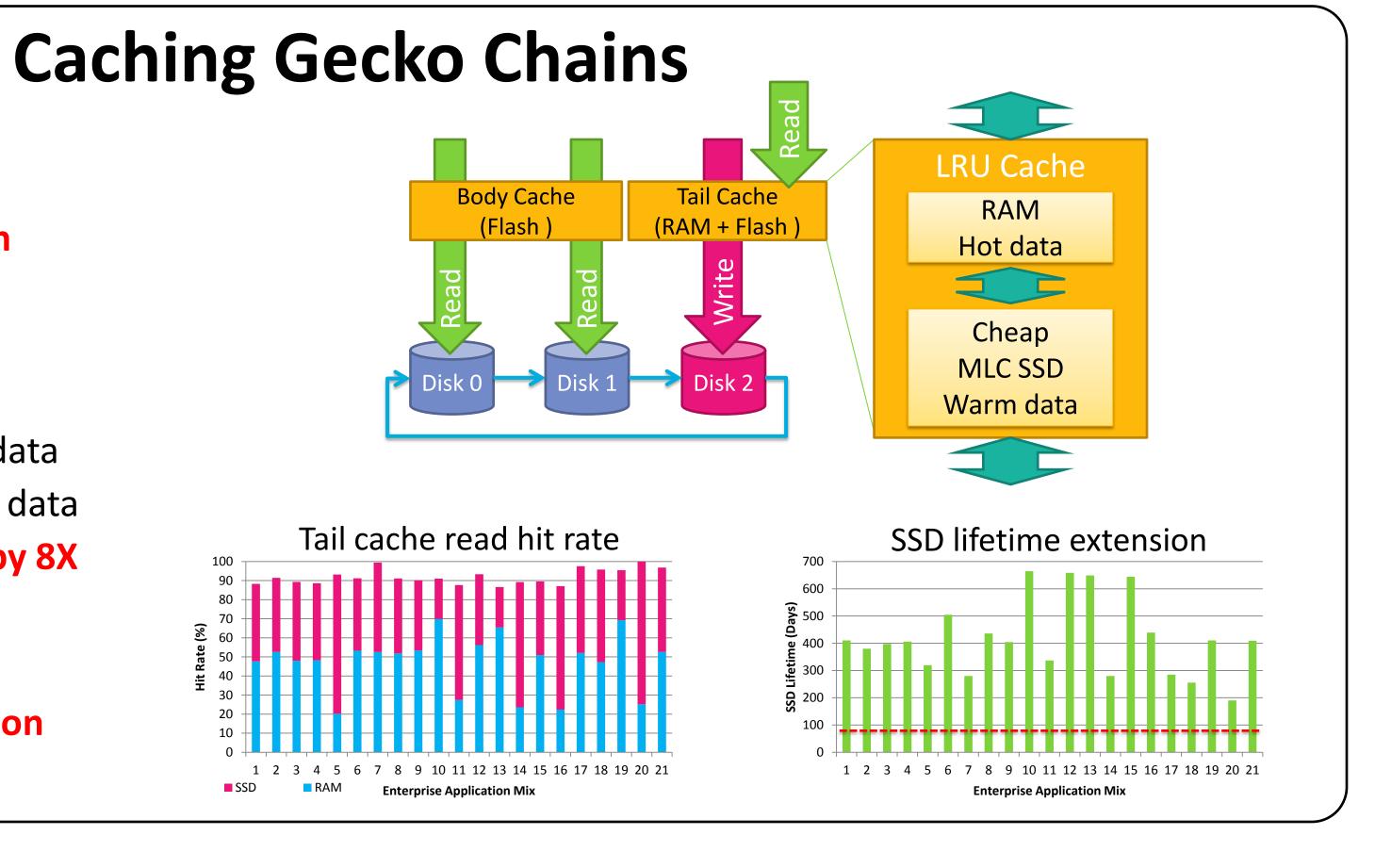
### Gecko: A Chain Logging Design

- Logs to one disk at a time to reduce I/O contention
- Logging (writing) takes place ONLY at the tail drive
- GC and read occurs at the rest of the drives
- Logging solves
  - Write-write contention Write-GC-write contention
- Chaining solves/reduces
- Write-GC-read contention
- Write-read contention Mirroring/Striping chains enables
- Power saving w/o consistency concerns
- High performance
- High fault tolerance



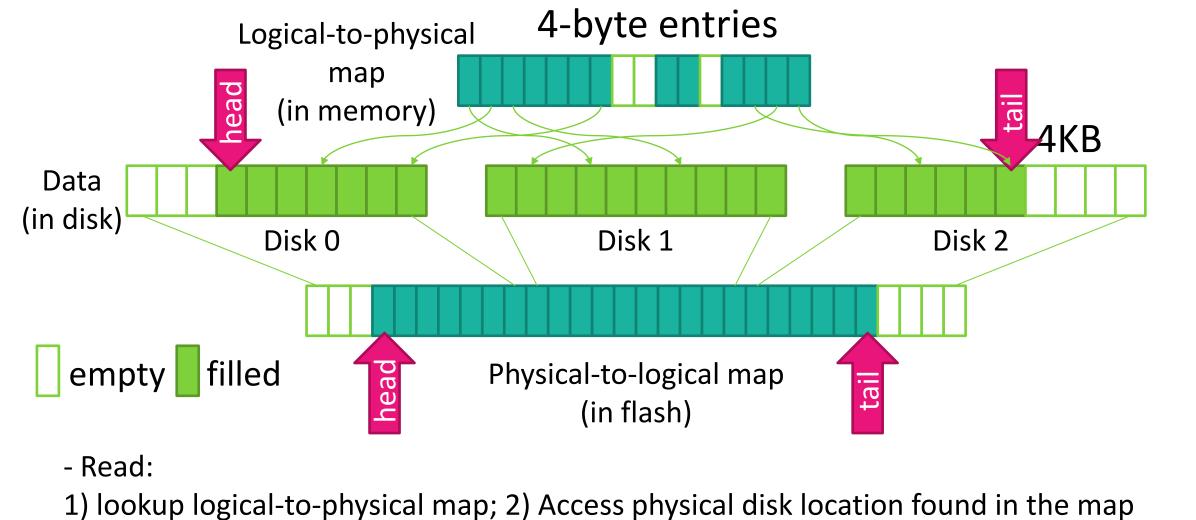
# Caching tail drive

- Reduces write-read contention
  - 86% of reads absorbed from mix of MS I/O traces
- RAM + SSD cache
  - RAM: Small amount of hot data
  - SSD: Large amount of warm data
- RAM extends SSD lifetime by 8X
- Caching body drives
- SSD only cache
- Reduces read-read contention
- Relatively low hit rate



# Metadata Management

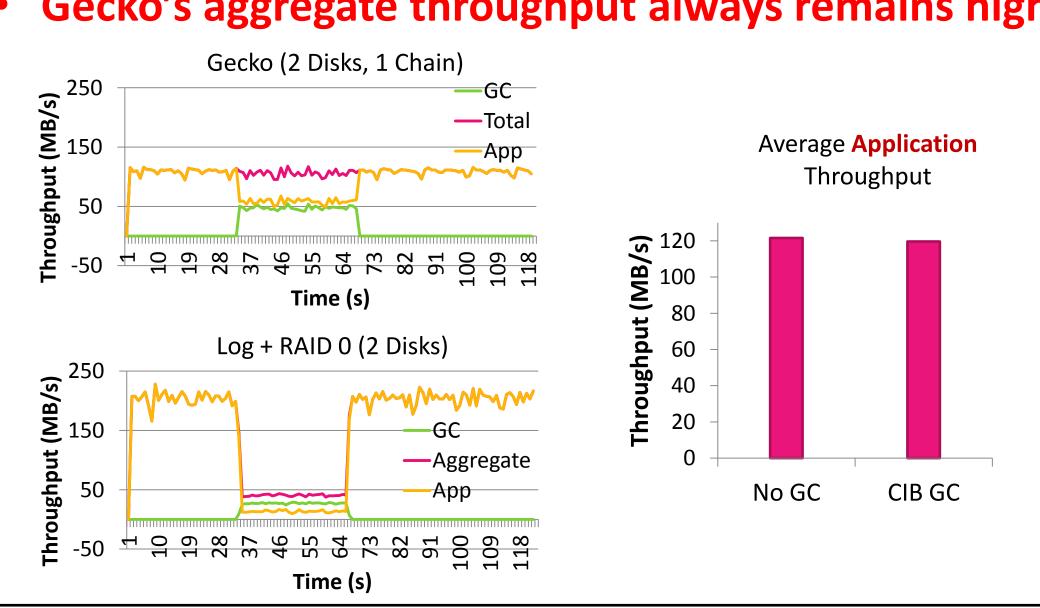
- In memory logical-to-physical map
  - 4-byte entries per page
  - 8GB for 8TB storage
- In flash physical-to-logical map
- Maintains persistence
- Flushed to flash every 1024 page writes
- Written in sequential order
  - High flash performance
- Good for flash lifetime



- 1) lookup logical-to-physical map; 2) Access physical disk location found in the map - Write:
- 1) Get next physical tail location X; 2) Update logical-to-physical map with X;
- 3) Maintain physical-to-logical map in memory
- 4) Flush in-memory physical-to-logical map to flash whenever 1024 entries are written - Recovery:
- 1) Scan physical-to-logical map in flash to reconstruct logical-to-physical map in memory

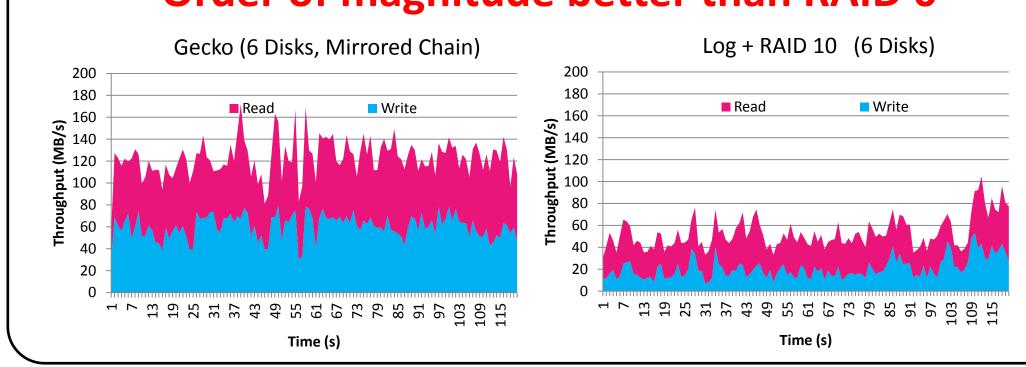
### Throughput under GC

- Write only synthetic workload on 2 disks
- Gecko's aggregate throughput always remains high



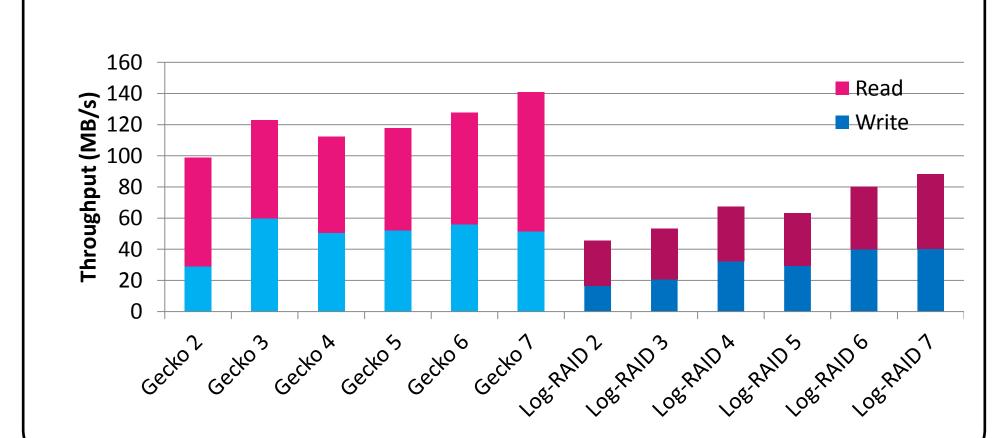
#### Running Enterprise Workloads

- Mix of 8 MSR Cambridge and MS enterprise workloads
- 6 Disk Configurations
- Gecko performs
- 2-3X better than Log+RAID10
- Order of magnitude better than RAID 0



#### Varying Chain Length

- SINGLE uncontended disk performs better than **SEVEN** contended disks
- Separating read and write reduces contention
- Typically 3 to 4 disk chains perform reasonably



## Conclusion

Higher

and reliable

throughput

- Gecko maintains high I/O performance
- Securing single uncontended disk
- Separating reads and writes
- Log-structured designs
- Oblivious to write-write contention
- Sensitive to GC-write and read-write contention
- Gecko fixes GC-write and read-write contention
- Log tail and head separation using chain logging
- Use of RAM+SSD tail disk cache